

# Machine-Level Programming V: Memory layout

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based on Jinyang Li's slides

# x86 Procedure Recap

## ■ **call**

- push return address on stack, jump to label

## ■ **ret**

- pop 8 bytes from stack into PC

## ■ **Argument passing from caller to callee**

- First 6 arguments passed in registers (%rdi, %rsi, %rdx, %rcx, %r8, %r9)
- Rest on stack

## ■ **Return value passing from callee to caller**

- %rax

## ■ **Local variable**

- either registers, or allocated on stack (deallocated before ret)

## ■ **Caller vs. callee-save registers**

- Caller-save: all “argument” registers, %rax, %r10, %r11
- Callee-save: %rbx, %r12, %r13, %r14, %rbp

# Recap: Procedure call example

```
int add2(int a, int b)
{
    return a + b;
}
```

```
int add3(int a, int b, int c)
{
    int r = add2(a, b);
    r = r + c;
    return r;
}
```

```
add2:
    leal    (%rdi,%rsi), %eax
    ret
```

```
add3:
    pushq  %rbx
    movl   %edx, %ebx
    movl   $0, %eax
    call  add2
    addl   %ebx, %eax
    popq  %rbx
    ret
```

a: %edi  
b: %esi  
c: %edx

%edx (containing c)  
is needed after call,  
so save in %ebx

## Registers

First 6 Arguments: %rdi, %rsi, %rdx, %rcx, %r8, %9

Return value: %rax

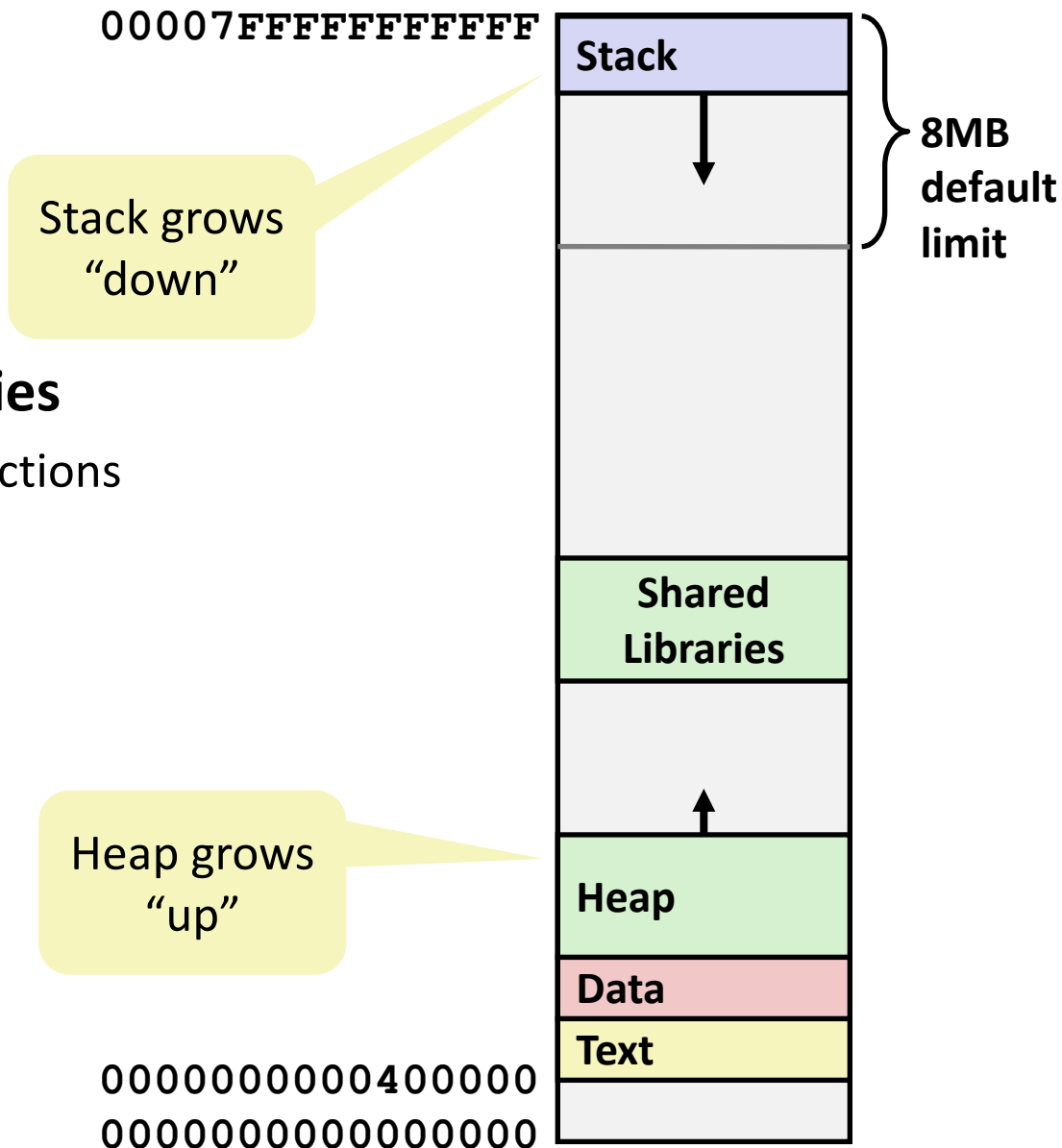
# OS loads a program to memory

- **OS loads different parts of a program into different memory regions**
- **Parts of a running program:**
  - Stack
    - e.g. local variables
  - Heap
    - e.g. malloc(), new
  - (statically allocated) Data
    - global variables, string constants
  - Executable instructions
- **Why different regions?**
  - need to grow independently
  - need different permissions

# x86-64 Linux Memory Layout

*not drawn to scale*

- Stack
- Heap
- Data
- Text / Shared Libraries
  - aka executable instructions



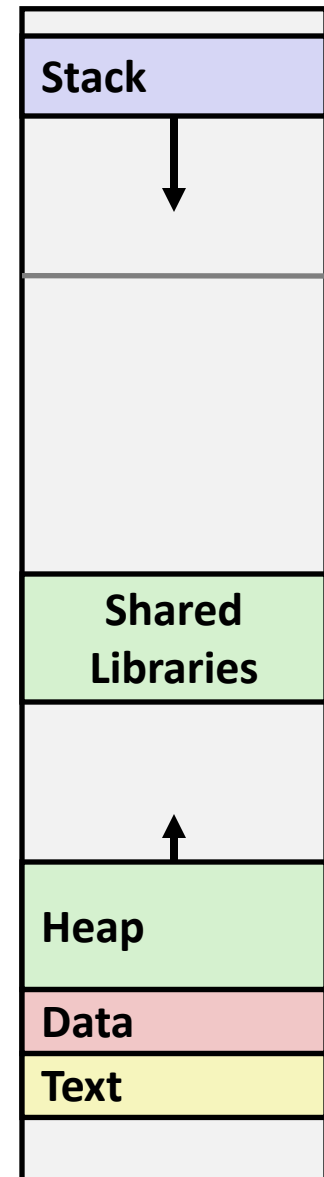
# Memory Allocation Example

```
char big_array[1<<24]; /* 16 MB */
char huge_array[1<<31]; /* 2 GB */

int global = 0;

int useless() { return 0; }

int main ()
{
    void *p1, *p2, *p3, *p4;
    int local = 0;
    p1 = malloc(1 << 28); /* 256 MB */
    p2 = malloc(1 << 8); /* 256 B */
    p3 = malloc(1 << 32); /* 4 GB */
    p4 = malloc(1 << 8); /* 256 B */
    ...
}
```

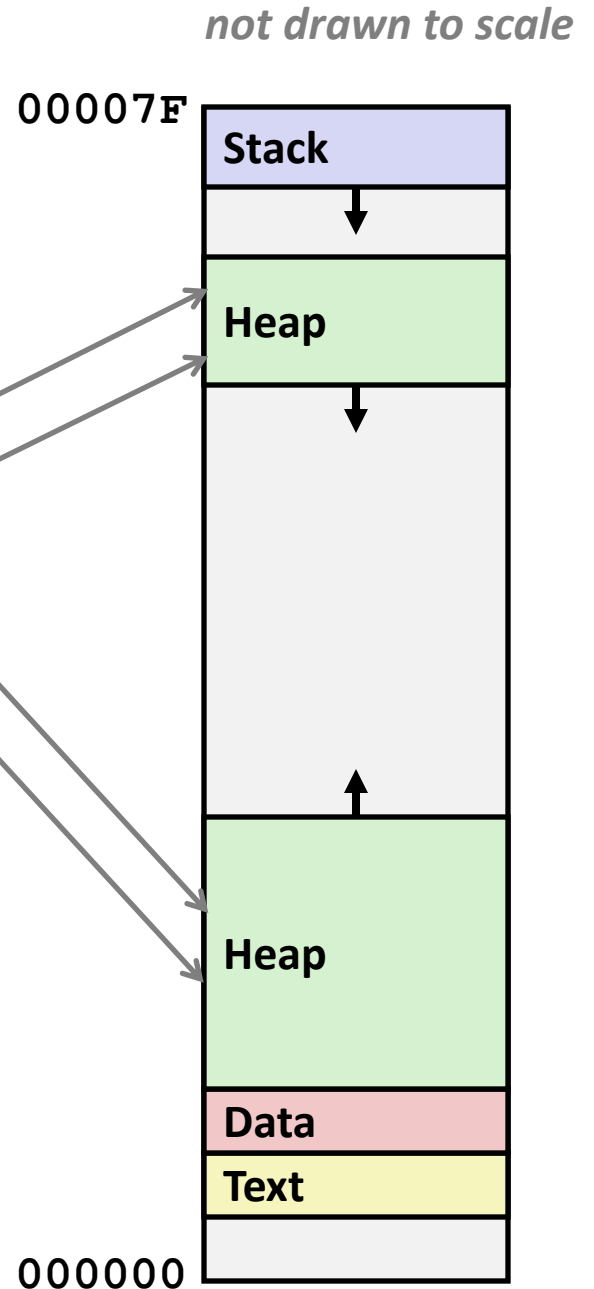


*Where does everything go?*

# x86-64 Example Addresses

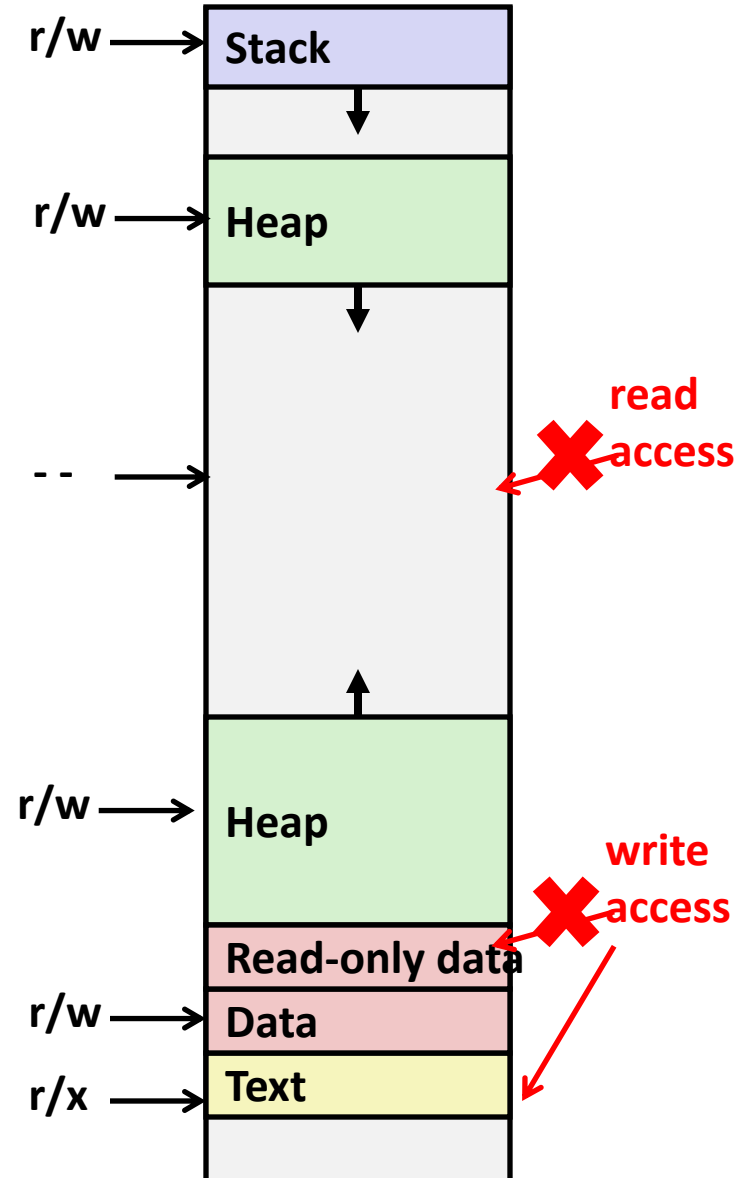
address range  $\sim 2^{47}$

local	0x00007ffe4d3be87c
p1	0x00007f7262a1e010
p3	0x00007f7162a1d010
p4	0x000000008359d120
p2	0x000000008359d010
big_array	0x0000000080601060
huge_array	0x0000000000601060
main()	0x000000000040060c
useless()	0x0000000000400590



# Segmentation Fault

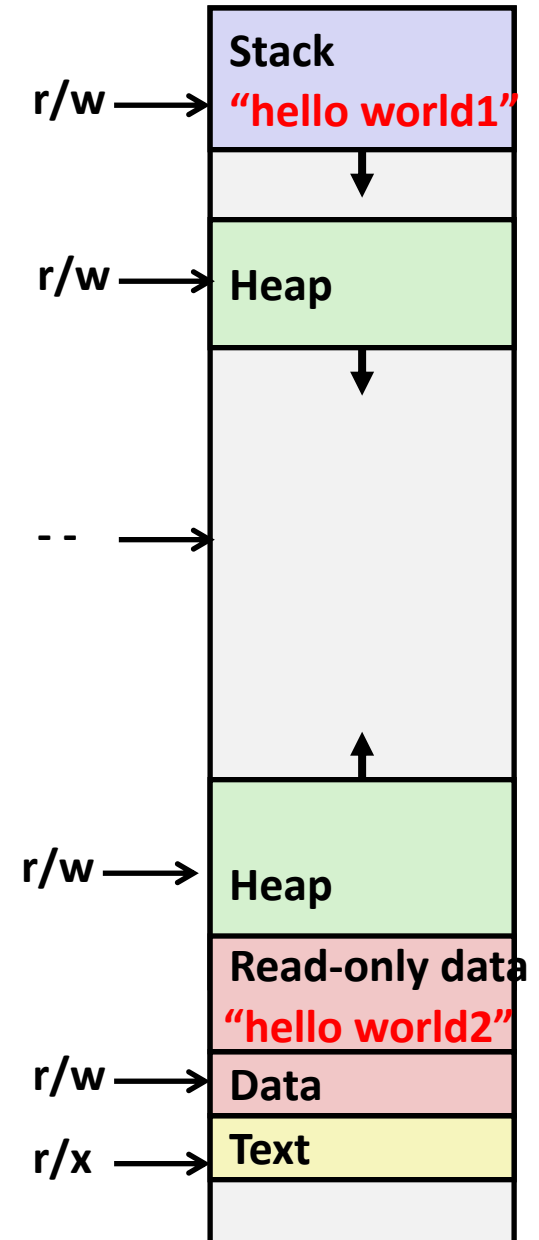
- Each memory segment can be readable (r), executable (x), writable(w), or none at all (-)
- Segmentation fault occurs when program tries to access “illegal” memory
  - Read from segment with no permission
  - Write to read-only segments





# Segmentation fault example

```
int main() {  
    char s1[100] = "hello world1";  
    char *s2 = "hello world2";  
    printf("str1 %p str2 %p\n", s1, s2);  
    s1[0] = 'H';  
    s2[0] = 'H';  
    ...  
}
```



# Not all Bad Memory Access lead to immediate segmentation

```
typedef struct {
    int a[2];
    double d;
} struct_t;

double fun(int i) {
    struct_t s;
    s.d = 3.14;
    s.a[i] = 1073741824; /* Possibly out of bounds */
    return s.d;
}
```

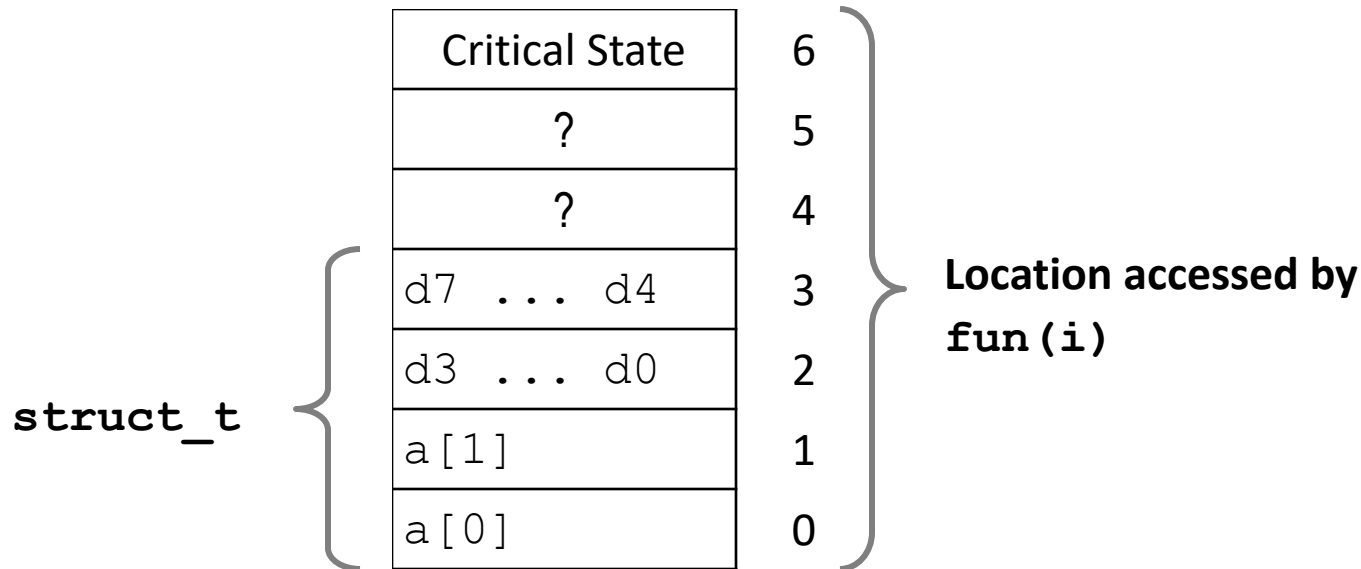
<b>fun (0)</b>	<b>→</b>	<b>3.14</b>
<b>fun (1)</b>	<b>→</b>	<b>3.14</b>
<b>fun (2)</b>	<b>→</b>	<b>3.1399998664856</b>
<b>fun (3)</b>	<b>→</b>	<b>2.00000061035156</b>
<b>fun (4)</b>	<b>→</b>	<b>3.14</b>
<b>fun (6)</b>	<b>→</b>	<b>Segmentation fault</b>

- Result is system specific

# Memory Referencing Bug Example

```
typedef struct {  
    int a[2];  
    double d;  
} struct_t;
```

fun(0)	→	3.14
fun(1)	→	3.14
fun(2)	→	3.1399998664856
fun(3)	→	2.00000061035156
fun(4)	→	3.14
fun(6)	→	Segmentation fault



# Such problems are a BIG deal

- **Generally called a “buffer overflow”**
  - when exceeding the memory size allocated for an array
- **Why a big deal?**
  - It's the #1 technical cause of security vulnerabilities
    - #1 overall cause is social engineering / user ignorance
- **Most common form**
  - Unchecked lengths on string inputs
  - Particularly for bounded character arrays on the stack
    - sometimes referred to as stack smashing